HEAPS

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Announcements

- PA03 checkpoint due tomorrow (05/31) at midnight. Submit to pa03-checkpoint assignment on gradescope
- PA03 due next Friday 06/08 at midnight. Submit to the pa03 assignment on gradescope
- Final exam on Wed 06/13 (8a -11a)
- Review session on Tuesday (06/05): first session 2:00p to 3:00p and the other from 3:00p to 4:00p

How is PA03 going?

- A. Done
- B. On track to finish
- C. Having trouble with the checkpoint (design)
- D. Just started
- E. Haven't started

Heaps: Supported Operations

Min-Heaps

BST

- Insert :
- Min:
- Delete Min:
- Max
- Delete Max

Choose heap if you are doing repeated insert/delete/(min OR max) operations

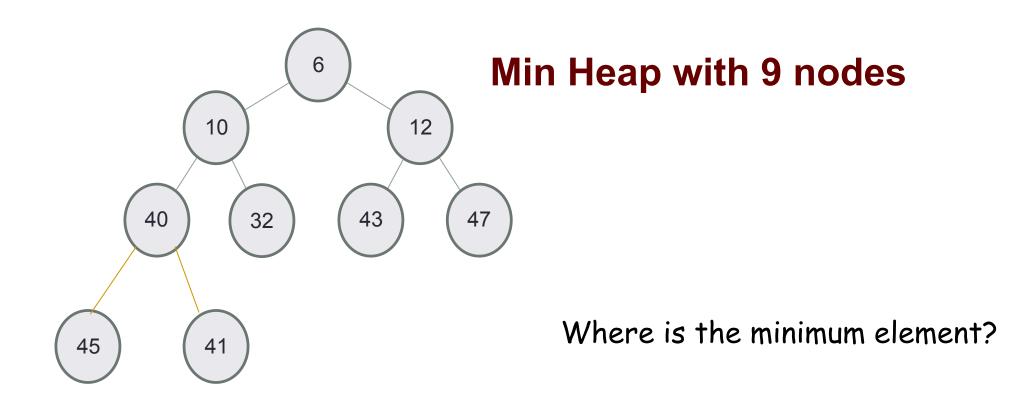
Max-Heap

Applications:

- Efficient sort
- Finding the median of a sequence of numbers
- Compression codes

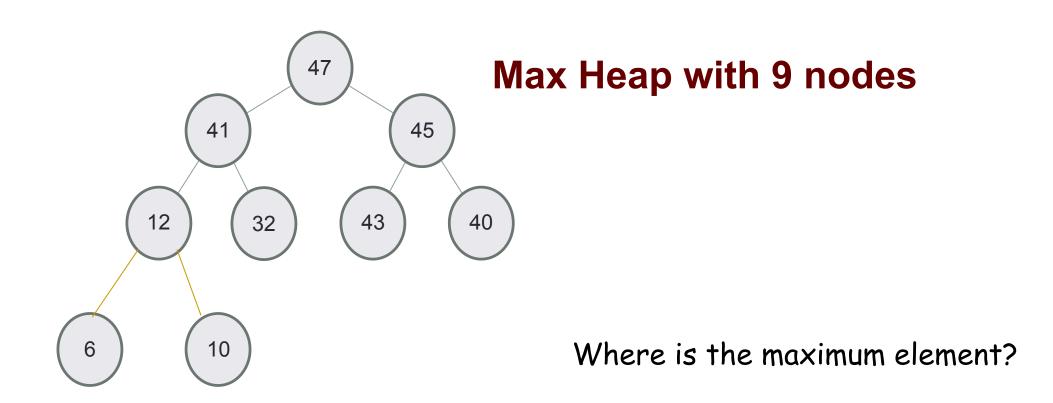
Heaps as binary trees

- Rooted binary tree that is as complete as possible
- In a min-Heap, each node satisfies the following heap property: key(x)<= key(children of x)



Heaps as binary trees

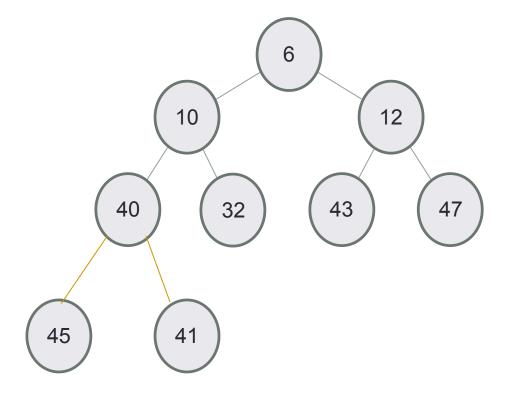
- Rooted binary tree that is as complete as possible
- In a max-Heap, each node satisfies the following heap property: key(x)>= key(children of x)



Identifying heaps

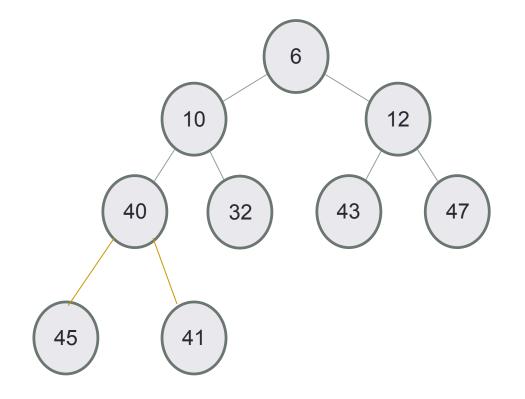
Starting with the following min Heap which of the following operations will result in something that is NOT a min Heap

- A. Swap the nodes 40 and 32
- B. Swap the nodes 32 and 43
- C. Swap the nodes 43 and 40
- D. Insert 50 as the left child of 45 E. C&D



Structure: Complete binary tree

A heap is a complete binary tree: Each level is as full as possible. Nodes on the bottom level are as far left as possible



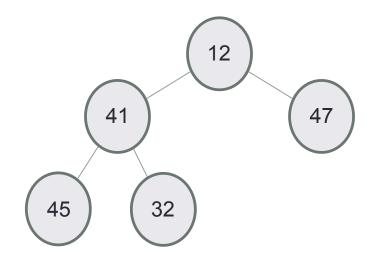
Insert into a heap

- Insert key(x) in the first open slot at the last level of tree (going from left to right)
- If the heap property is not violated Done

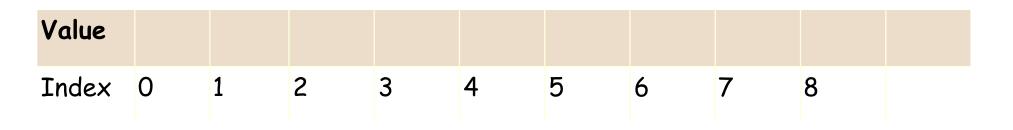
Insert the elements {12, 41, 47, 45, 32} in a min-Heap

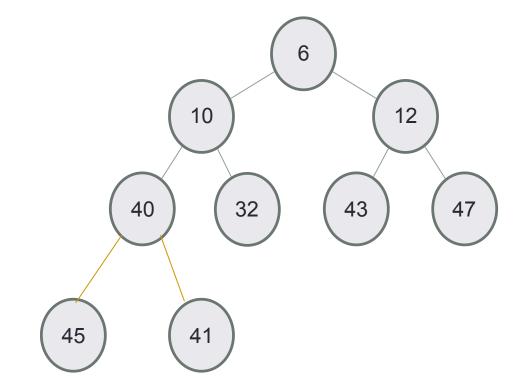
Insert 32 into a heap

- Insert key(x) in the first open slot at the last level of tree (going from left to right)
- If the heap property is not violated Done
- Else: while(key(parent(x))>key(x)) swap the key(x) with key(parent(x))



Implementing heaps as arrays



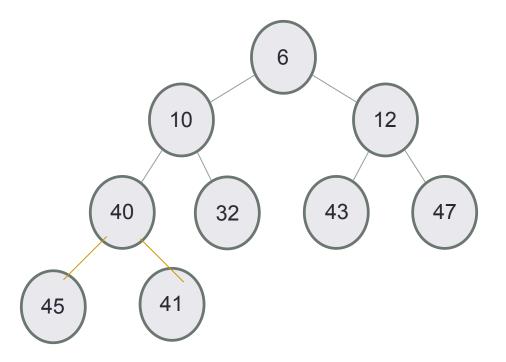


How is the array implementation of the heap useful?

- More space efficient
- Accessing parent and children of a node is O(1)
- Easier to insert elements in the heap

Conceptualize heaps as trees, implement as arrays

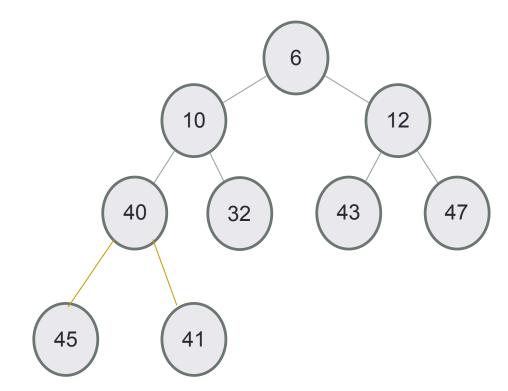
Value	6	10	12	40	32	43	47	45	41	
Index	0	1	2	3	4	5	6	7	8	



For a node at index i, what is the index of the left and right children?

- A. (2*i, 2*i+1)
- B. (2*i+1, 2*i+2)
- C. (log(i), log(i)+1)
- **D.** None of the above

Conceptualize heaps as trees, implement as arrays



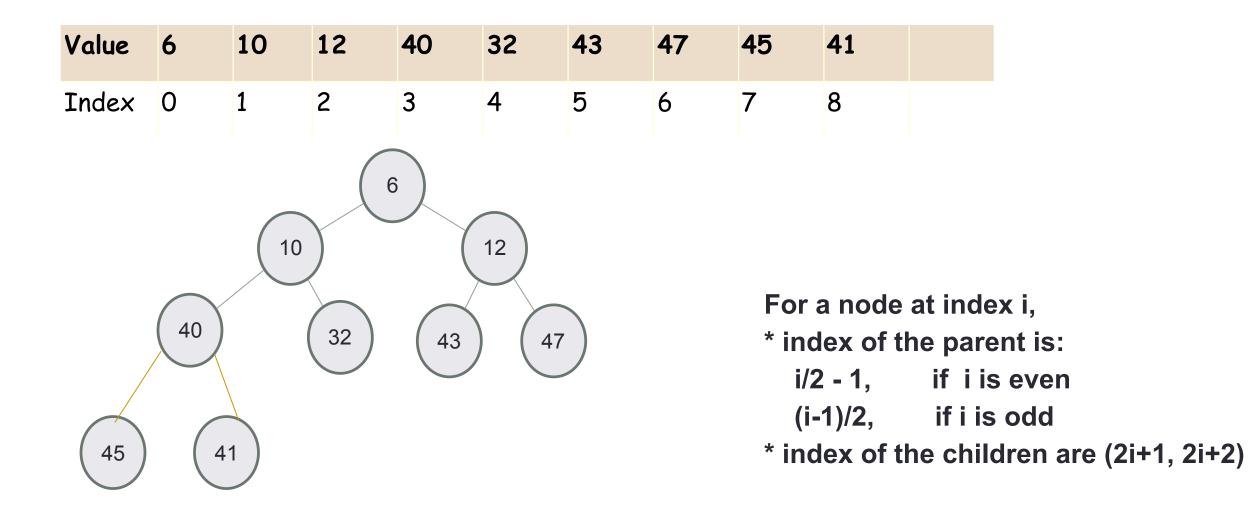
For a node at index i, index of the parent is:

i/2 - 1,	if i	is	even
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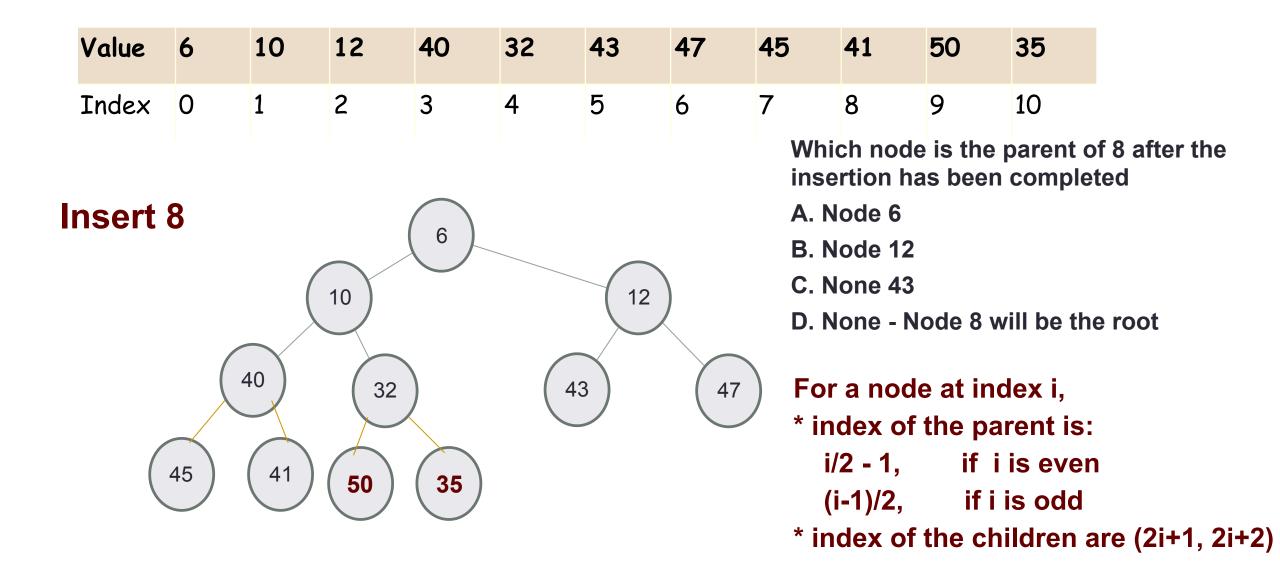
(i-1)/2, if I is odd

Value	Index	Index of parent	Index of children
6	0	-	1, 2
10	1	0	3, 4
12	2	0	5,6
40	3	1	7,8
32	4	1	
43	5	2	
47	6	2	
45	7	3	
41	8	3	

Insert 50, then 35

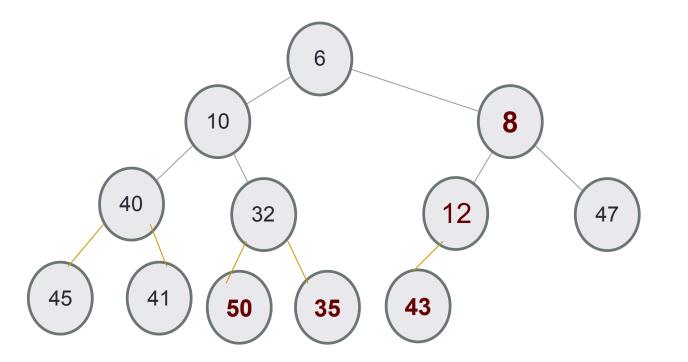


Insert 8 into a heap



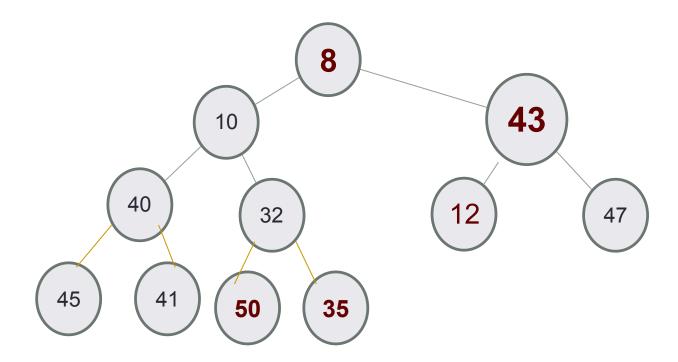
Delete min

- Replace the root with the rightmost node at the last level
- "Bubble down"- swap node with one of the children until the heap property is restored



Delete min

- Delete the root:
 - Replace the root with the last node in the array
 - If heap property is violated swap with the child that has the LOWEST key value, repeat until heap property is restored



To fix the heap property on a delete BUBBLE DOWN! Worst case: O(logN)

Applications

- Efficient sort
- Finding the median of a sequence of numbers
- Compression codes